



# INTERNATIONAL JOURNAL OF PURE AND APPLIED RESEARCH IN ENGINEERING AND TECHNOLOGY

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## AN APPROACH TO WIRELESS SECURITY BY HANDLING COMPUTATIONS

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Accepted Date: 27/02/2014 ; Published Date: 01/05/2014

**Abstract:** Research on security in wireless has been an interesting active area since the last decade. Generally cryptographic tools use is more in wireless security studies, but these tools based on computational assumptions, which may turn out to be invalid in the future. so there is need to build cryptographic tools which does not depend on computational assumptions. 1-out-of-2 oblivious transfer tool plays a major role in cryptography & can build a cryptographic protocol for any polynomial-time computable function which does not depend on any computational assumption. private communications and the privacy preserving password verification are some application of this protocol. This protocol fully implemented on wireless devices and conducted experiments in real environments.

**Keywords:** Wireless Security, Research, Handling Computation



PAPER-QR CODE

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How to Cite This Article:

Sachin Chawhan, IJPRET, 2014; Volume 2 (9): 1152-1157

## INTRODUCTION

Research on security in wireless domain has been an interesting active area since the last decade. Generally to provide the security in wireless area an encryption, authentication & the key agreement are essential. These cryptographic tools are really good for use but a common weakness they have, they are based on computational assumptions. For example, symmetric key encryption. The other encryption schemes, e.g., AES. However, when we use AES to encrypt a message, it can be an implicit assumption: the AES block cipher is a pseudorandom permutation. Due to the above assumption of pseudo randomness which is based on the cryptologists understanding of the *current* attacks on encryption schemes. In the future it may be possible, the AES scheme will be broken by newly invented cryptanalysis techniques. Some of the famous hash functions, including MD5 and SHA-0 work as a collision-resistant but cryptologists found that these assumptions are invalid. It will be necessary if we can remove cryptographic tools' dependence on such computational assumptions. It's a challenging problem to remove the computational assumptions from these tools. Here we are going to use wireless channel characteristics which used to build the cryptographic tool. These wireless channel characteristics are used to develop a tool known as 1-out-of-2 oblivious transfer. For simplicity it is known as OT(Oblivious transfer). using OT one can use cryptographic protocol for any polynomial time computable function. The other protocols that can build using OT are electronic voting protocol, digital cash protocols etc. So using a this protocol we can build other protocols which will be independent of computational assumptions. Both communicating devices share a Common secret key to communicate with each other privately. It can be observed that the security of such type of protocol based depend on the wireless channel characteristics, not on computational assumptions. The OT can be use as a protocol for privacy preserving password verification. In this method first communicating wireless device can verify a password from another communicating wireless device in such a way that the password is not known to either the first communicating wireless device nor by any attacker. In this paper the following major points highlighted.

- The main objective is to develop an OT protocol which is depend on the physical characteristics. This protocol does not depend on any computational assumption. It is a crucial step to build a strong wireless system security & without using computational assumptions.
- By using this protocol we can provide a private communications and a technique of privacy preserving password verification. related work

This tool is basically used to construct complex cryptographic protocols. Rabin was the first who proposed this tool first time. After some other people like Lempel, Even, Goldreich has propose

OT, an important variant of OT. The main advantage of OT is its completeness. Any two party protocol can be built by using OT shown by Kilian and due to this result it can be more beneficial by using it in case to construct multiparty protocols. Kilian and Crépeau present OT protocol which is based on noisy channels. To increase the efficiency Crépeau also proposed another OT protocol. Here the noisy channels they used are simple discrete memory less channels. In this approach we will use OT protocol which is based on wireless channels characteristics and which has more efficiency. We have demonstrated two applications of our OT21 protocol, private communications and privacy preserving password verification. In fact, there have been a number of works on private communications using the secrecy capacities of the wireless channel, among others. In particular, Vasudevantry to defend against the eavesdroppers by sending artificial noises to them, and focus on the scaling laws of secret communications without computational assumptions. In contrast, our private communications method is more practical in the sense that it does not need to control the received signals at the eavesdroppers. On the other hand, we stress that our private communications method is to illustrate the application of our OT21 protocol. We choose this application because it is simple and easy to understand, *not* because our private communications method is more efficient than the existing works on private communications.

### III. MODEL OF WIRELESS CHANNEL:

Suppose we have two parties  $A$  and  $B$  connected through a wireless channel between them. The magnitude of in phase component should be  $h$ , which follows a Gaussian distribution. Here  $h$  can be considered as a stochastic process; to represent the value of  $h$  at time  $t$  we use  $h(t)$ . One thing is noted that both the parties do not know the precise values of  $h(t)$ . Only they have to make an estimate. To initiate the communication process there will be a sharing of probe signal  $s(t)$ . Suppose at time  $t_1$ ,  $A$  receives a probe signal sent by  $B$ . Now party  $A$  sends a probe signal and party  $B$  receives it at time  $t_2$ . Both the parties now estimate the channel using their received probe signals.

The received signals of  $A$  and  $B$  can be shown as follows

$$r_a(t_1) = h(t_1)s(t_1) + n_a(t_1) \dots \dots \dots (1)$$

$$r_b(t_2) = h(t_2)s(t_2) + n_b(t_2) \dots \dots \dots (2)$$

where  $n_a(t_1)$  shows a receiver noise at party  $A$  and  $n_b(t_2)$  shows a receiver noise at party  $B$ .

### IV. APPLICATION : PRIVATE COMMUNICATIONS

This application deals with the private communication between two communicating parties by using a common secret key  $k$  between them. This is similar to the symmetric key encryption and

also used in decryption method of cryptography. The requirements for this communication are initially  $A$  send confidential message and  $B$  reply for the same, the condition is that the key should be same. If the key is different then there will be no reply from  $B$  when  $A$  send a confidential message.

The message send by  $A$  does not known to any eavesdropper. This method of private communication is very much similar to the symmetric key encryption and also to the symmetric key decryption used in cryptography. But this method is not identical to symmetric key encryption and the decryption. This model is different than the cryptography model we generally used, since here in this model of private communication we are not using ciphertext in traditional sense.

#### V. APPLICATION : PRIVACY PRESERVING PASSWORD VERIFICATION

To identify the user password verification is one of the method which is popular today. In wireless LANs, Many base stations at the beginning of sessions authenticate users using their passwords. But the risk of hacking of password is more when users send passwords by using wireless links. So it is necessary to provide privacy protection of password ,as password is going to use for authentication. In this method one wireless device has to identify the password from other device and with privacy preserving of the password. Following are the major requirements for this method. Initially wireless device  $A$  has to send a password entity that matches by the  $B$ 's record and suppose it should be matches by it then only accepted. Another possibility is that if it is not matches by the record of  $B$  then it will be rejected by  $B$ . The thing should note here is only  $A$  has to match a password with  $B$ 's record and nothing more than it. The same thing is happened opposite in case for the  $B$  also which learn nothing about the password rather it matches or not only. It is an advantage of this application is that the password matching process and off course the password does not known to the eavesdroppers that can attack on this process which is not possible due to privacy preserving password verification technique.

#### VI. CONCLUSION

Here , we present an OT protocol with two applications. This protocol does not depend on computational assumption and it is a major advantage of this protocol. It is found that in previously the major cryptographic tools also broken by the attackers and hence the need for such a technique which can able to provide strong security in wireless area is the requirement which can fulfill through this approach. Our work considered to be a step towards building wireless security which do not depends on computational assumptions although if we consider

the speed in term of performance then this protocol is still not that much fast as the generally used OT based on computational assumptions.

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